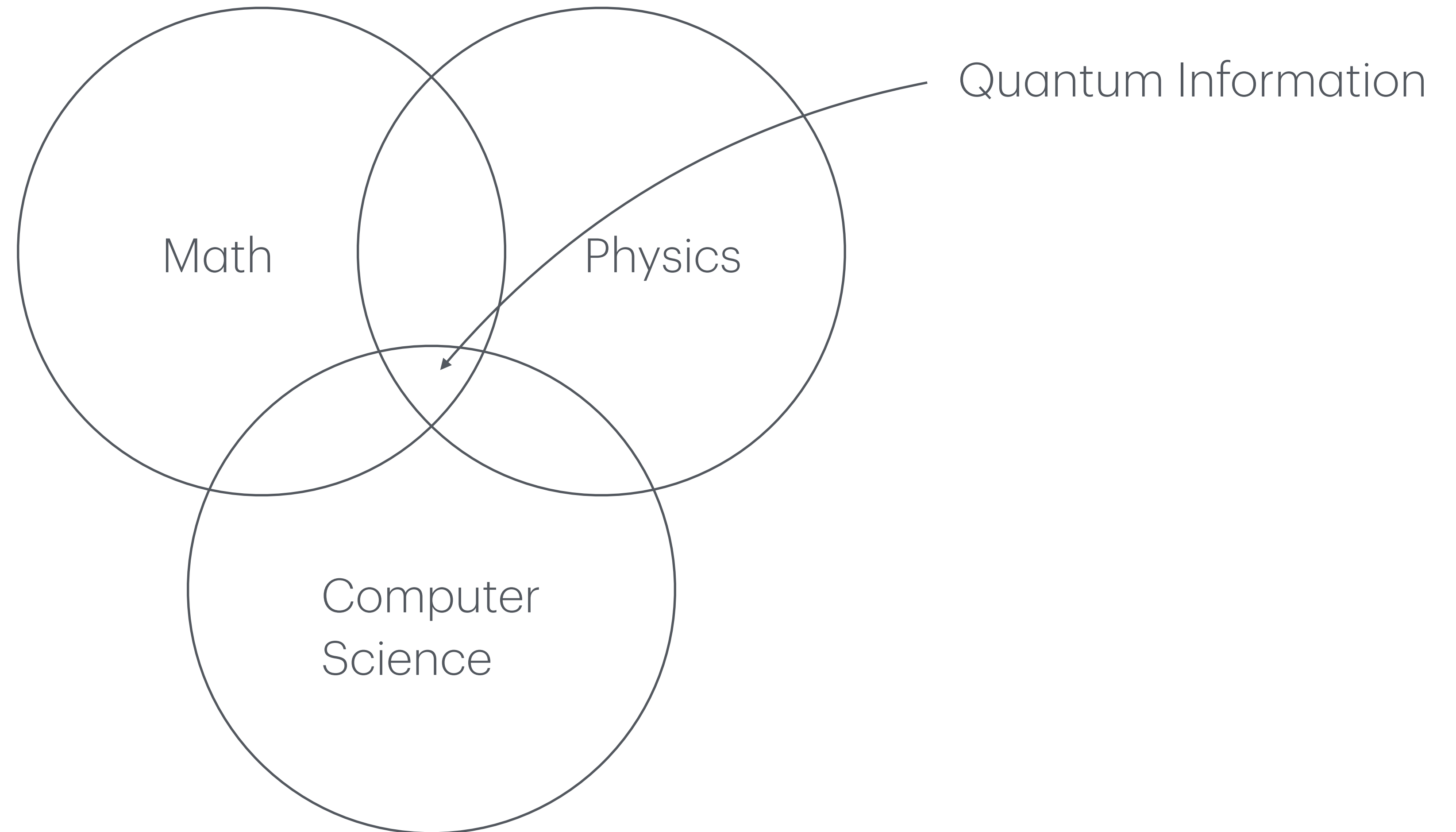
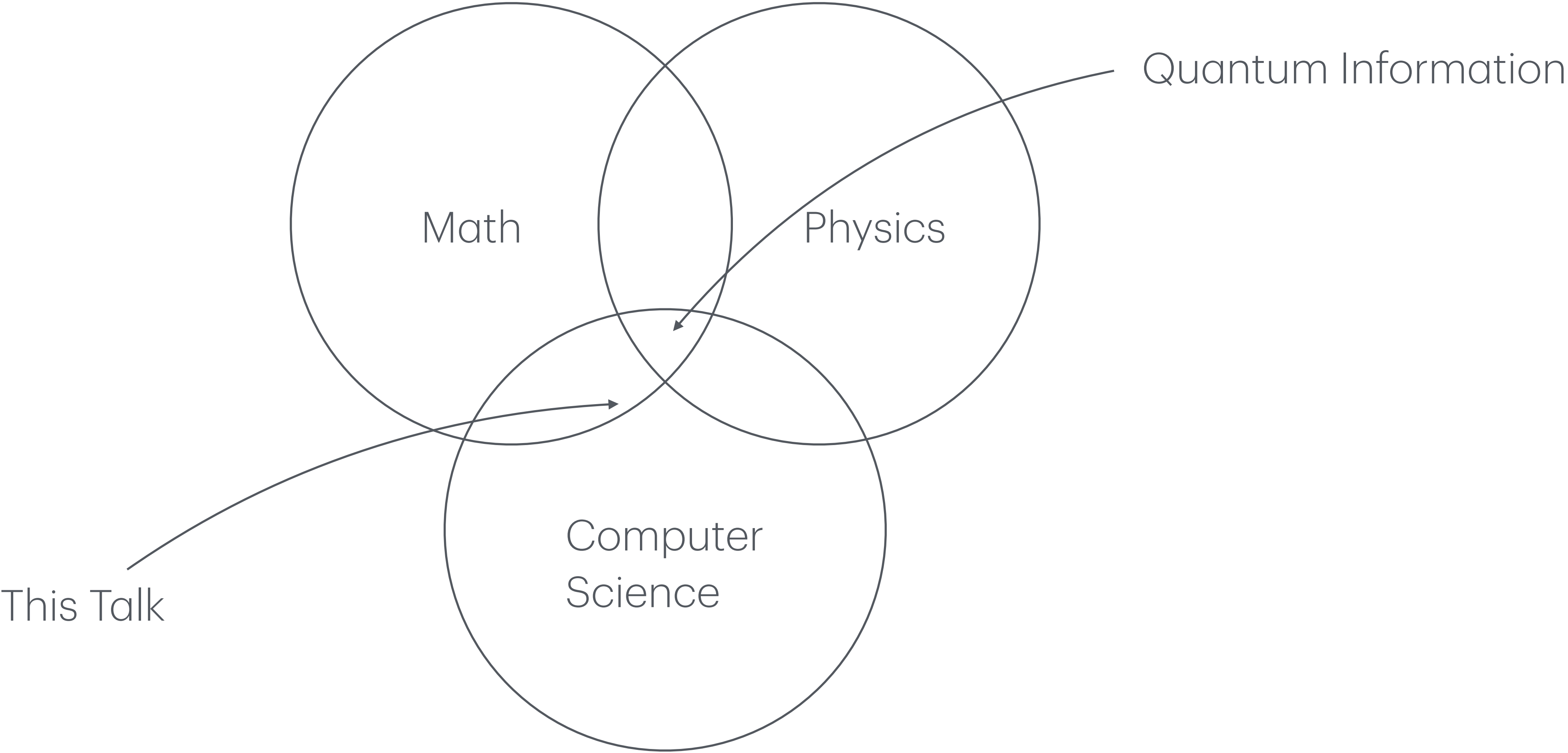


# Introduction to Quantum Information

# Quantum Information



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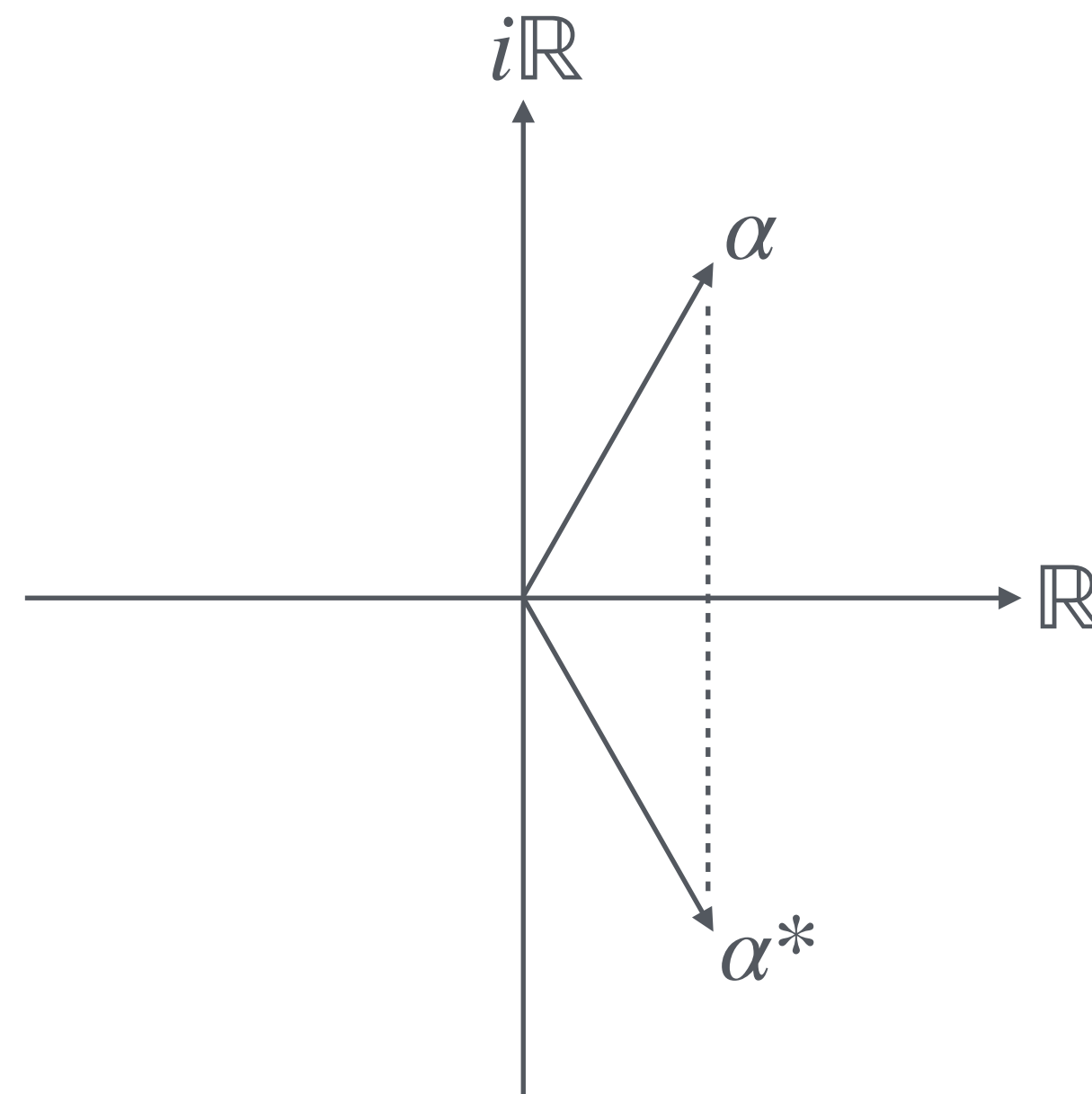
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- Of course nothing stops us from introducing more qubits to the system...

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- Any (pure) quantum state can be written as:

$$|\psi\rangle = \sum_x \alpha_x |x\rangle \quad \sum_x |\alpha_x|^2 = 1$$

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# Manipulating Quantum States

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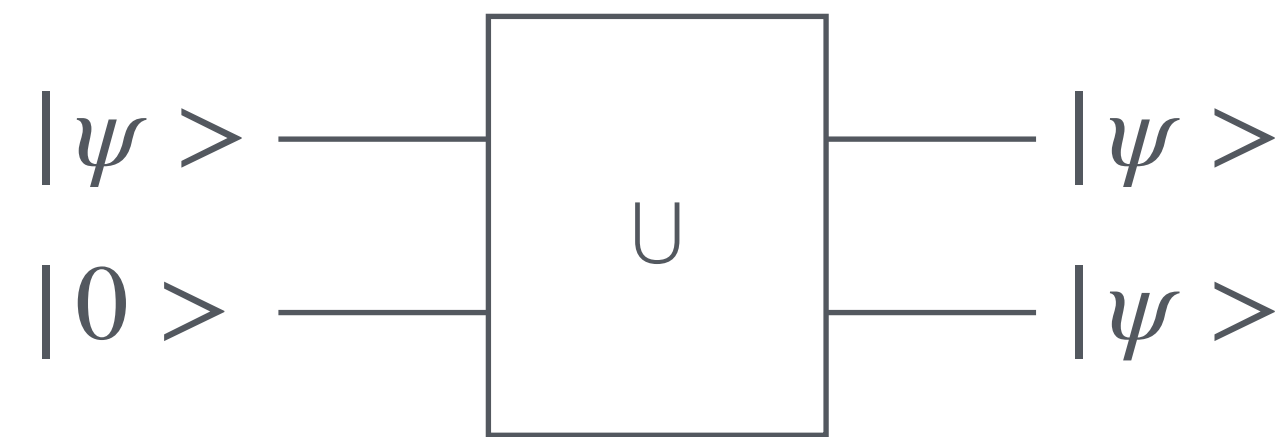
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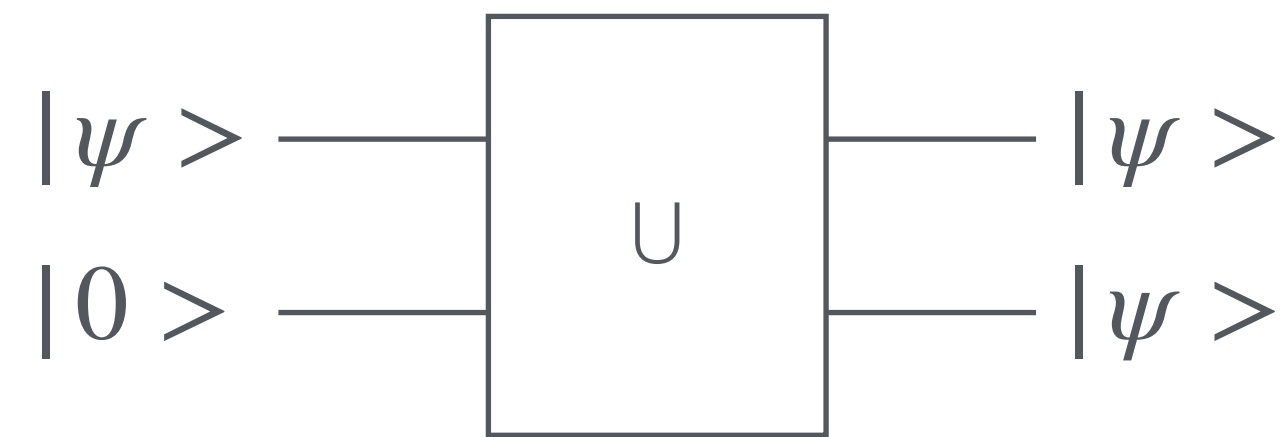
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- This implements the mapping

$$\begin{aligned} U(\alpha|0\rangle + \beta|1\rangle) \otimes |0\rangle &= (\alpha|0\rangle + \beta|1\rangle) \otimes (\alpha|0\rangle + \beta|1\rangle) \\ &= \alpha^2|00\rangle + \alpha\beta|01\rangle + \alpha\beta|10\rangle + \beta^2|11\rangle \end{aligned}$$

which is not linear, and in particular not unitary

# Computing with Quantum States

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- Now that we know the rules, let us see how/what we can compute in this model



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- So the question is, what kind of functions admit a “reversible” implementation?

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- THEOREM: If  $f$  can be implemented using  $s$ -many NAND gates, it can be implemented using  $O(s)$ -many Toffoli gates

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- In complexity theory terms:  $P \subseteq BQP$

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- The question to ask is: Are there quantum computations that we cannot simulate *efficiently* with classical computers?

# Quantum Circuits

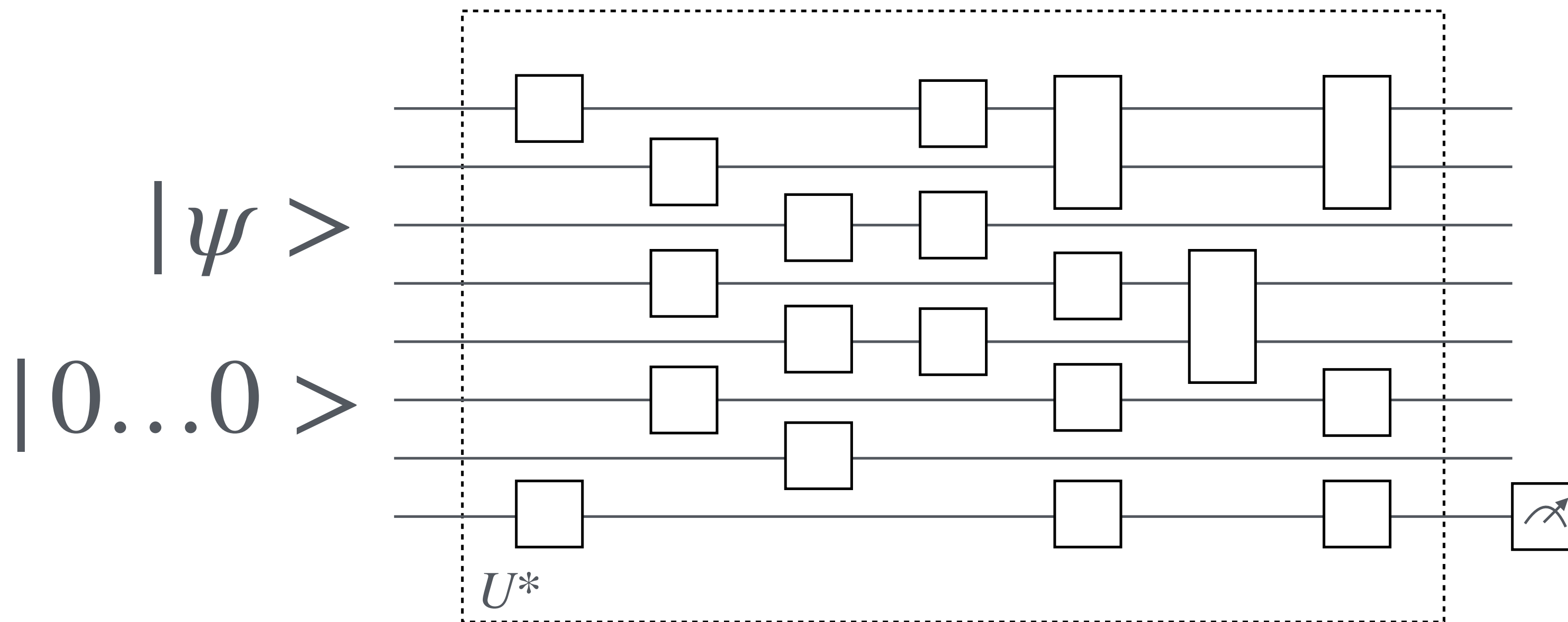


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$$T = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$$

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- The Solovay-Kitaev theorem bounds the number of gates (for *any* universal gate set) needed to approximate any (constant-dimension) unitary up to  $\epsilon$  precision, by  $\text{poly-log}(1/\epsilon)$

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- The more general q-ary QFT is also efficiently computable, where  $\omega_q = e^{2\pi i/q}$  is the q-th root of unity

$$QFT_q |x\rangle = \frac{1}{\sqrt{N}} \sum_y \omega_q^{x \cdot y} |y\rangle$$

Bonus: The CHSH Game

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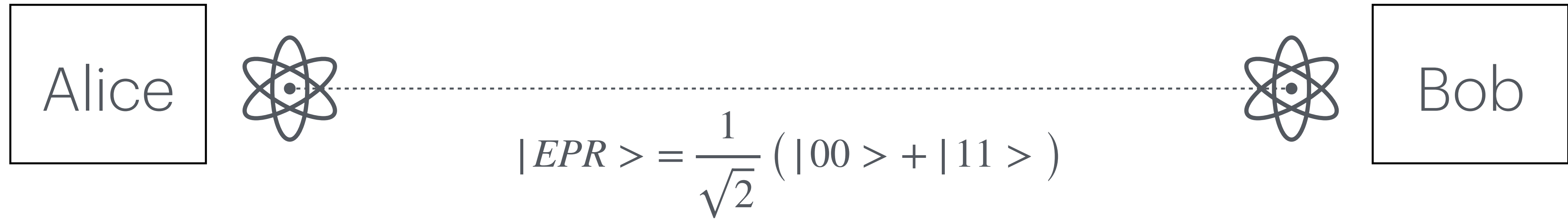
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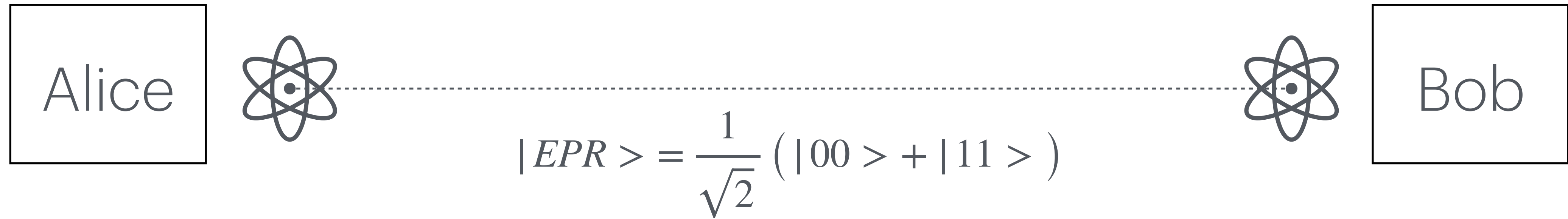


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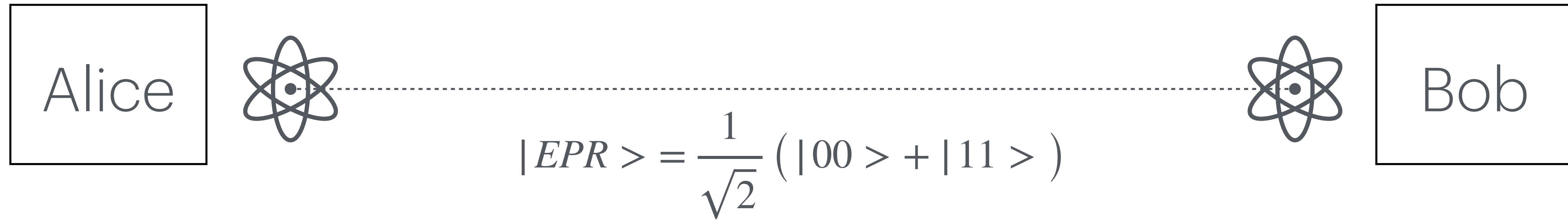


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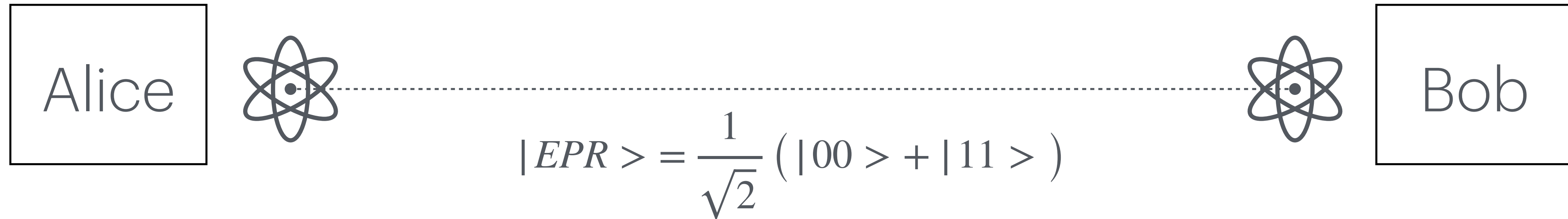
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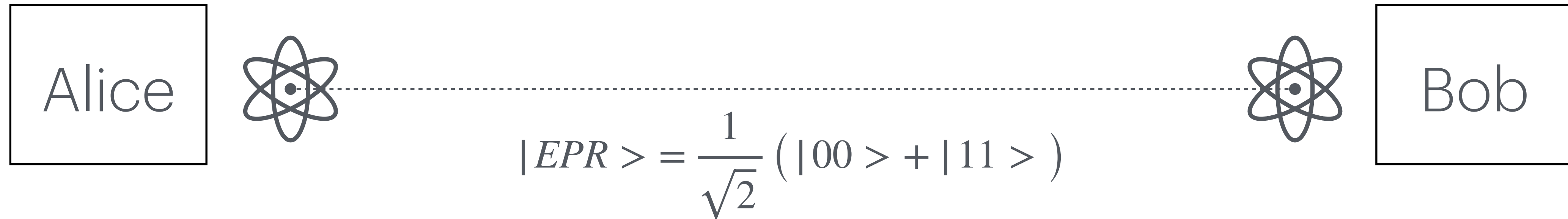
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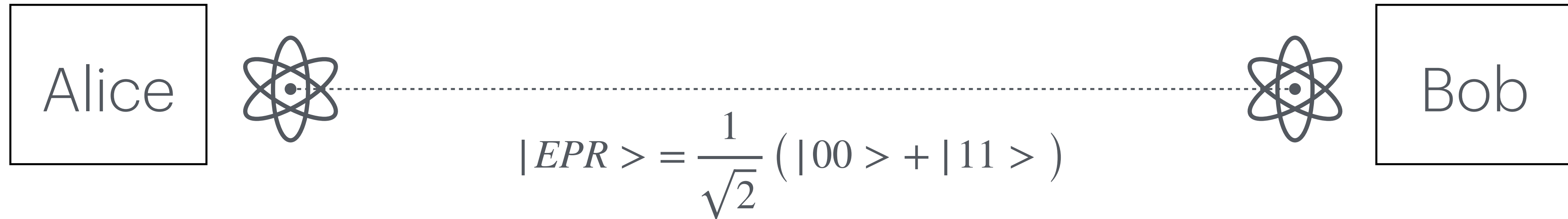
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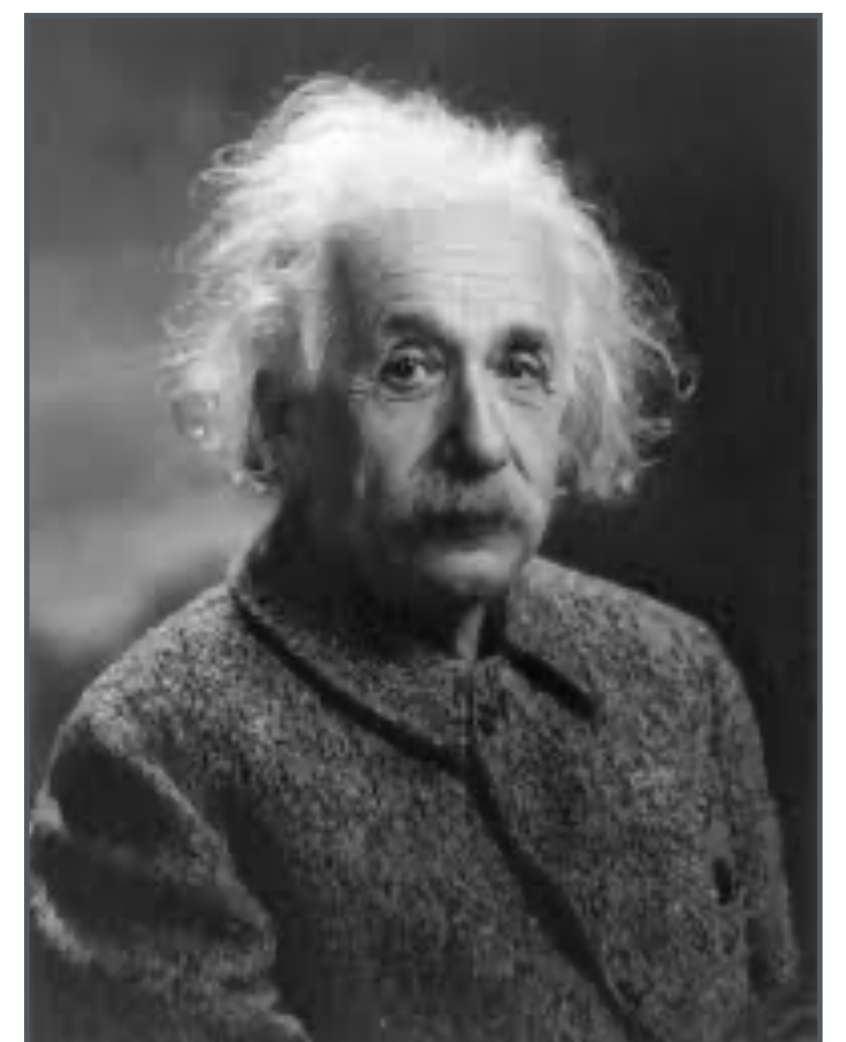
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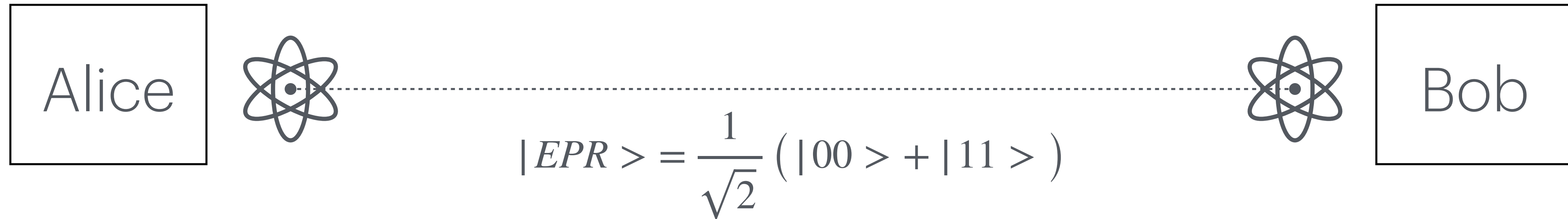
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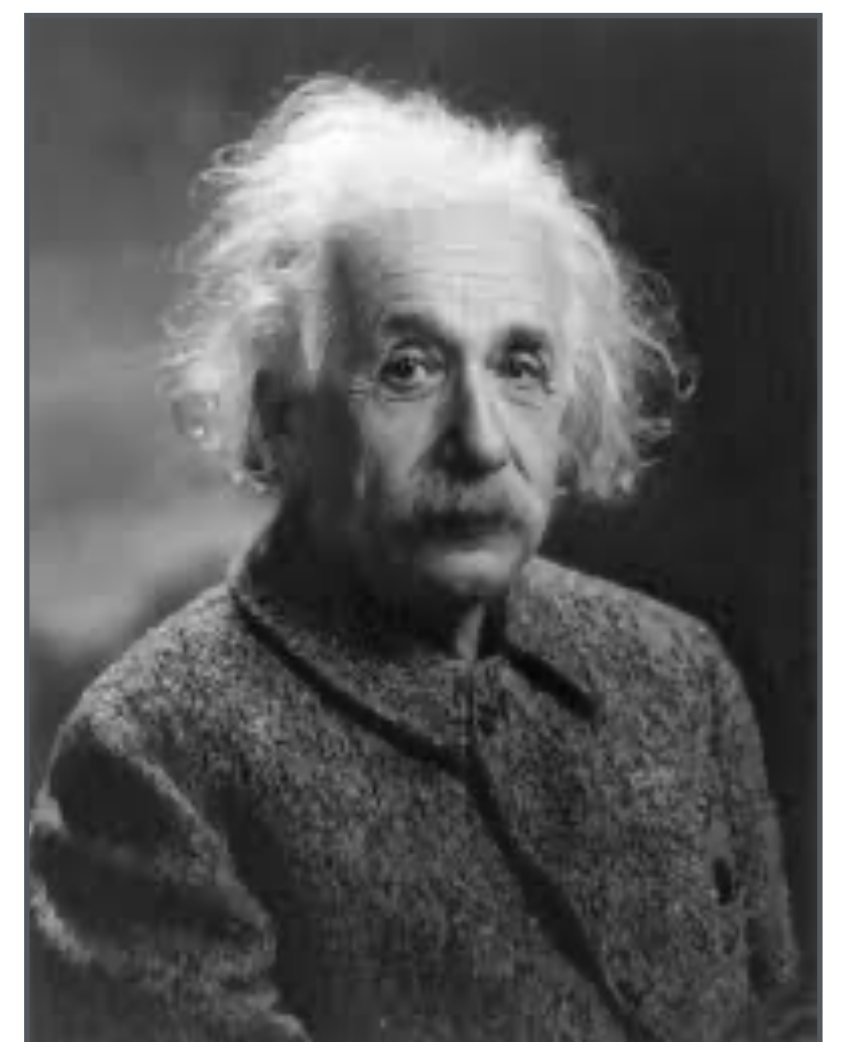


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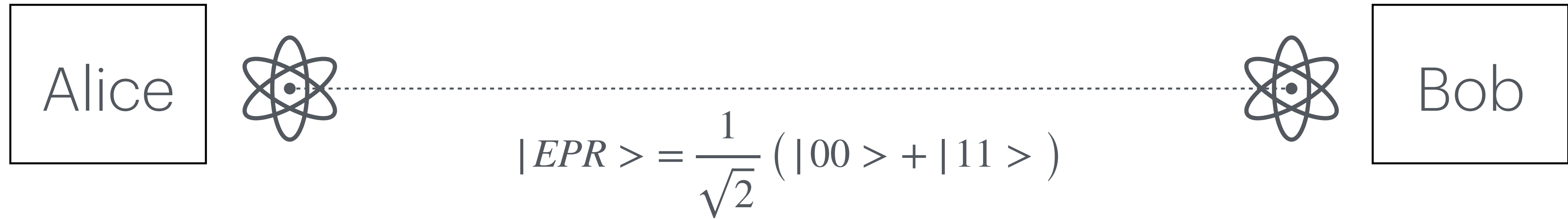


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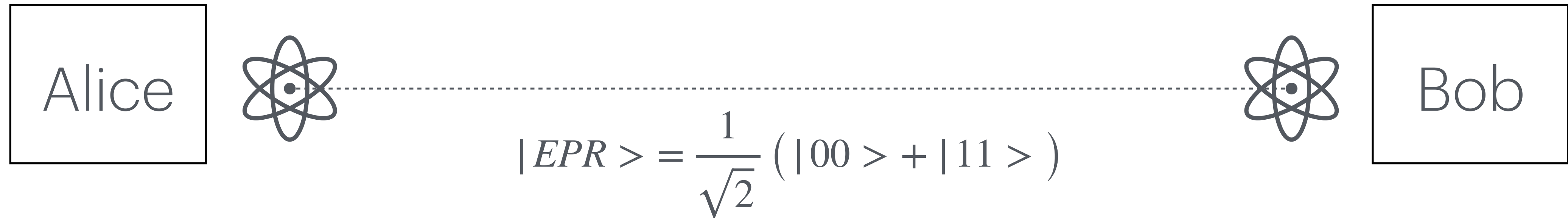
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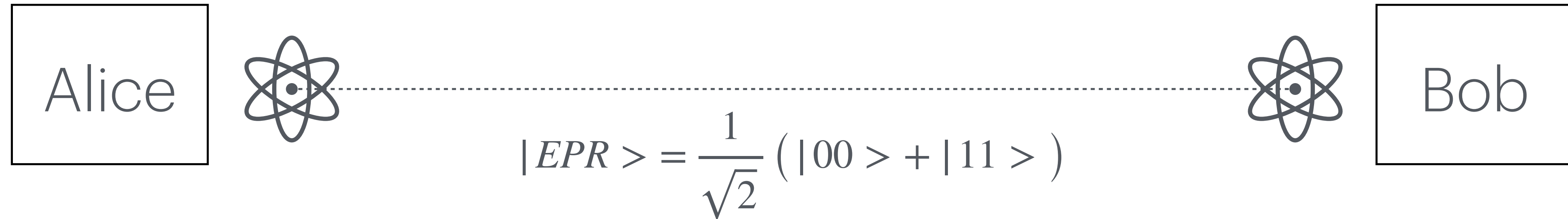
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Quantum Mechanics

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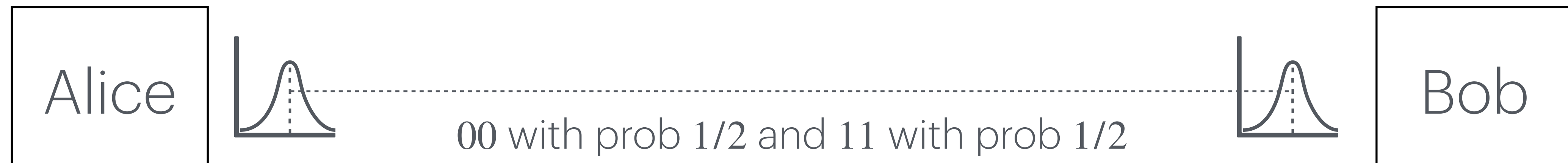
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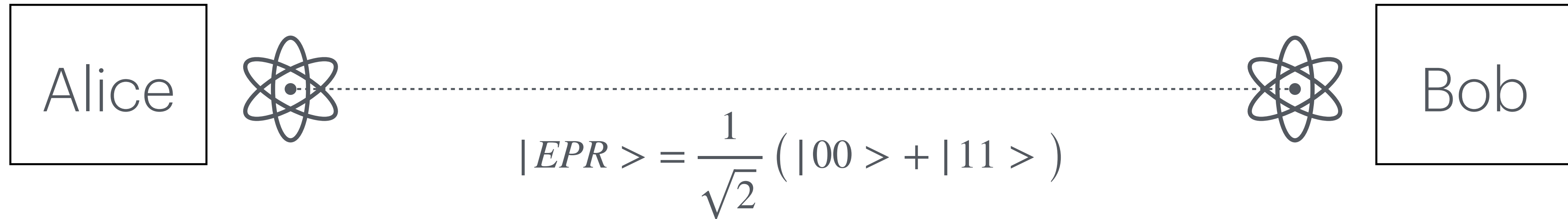
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Local Hidden Variables

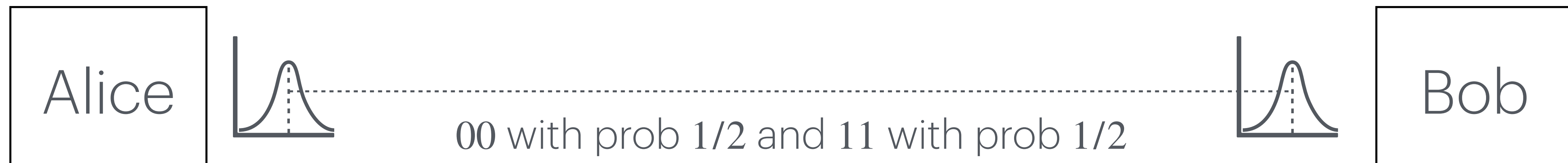


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Quantum Mechanics

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- PROBLEM: The probabilities are identical in both cases!

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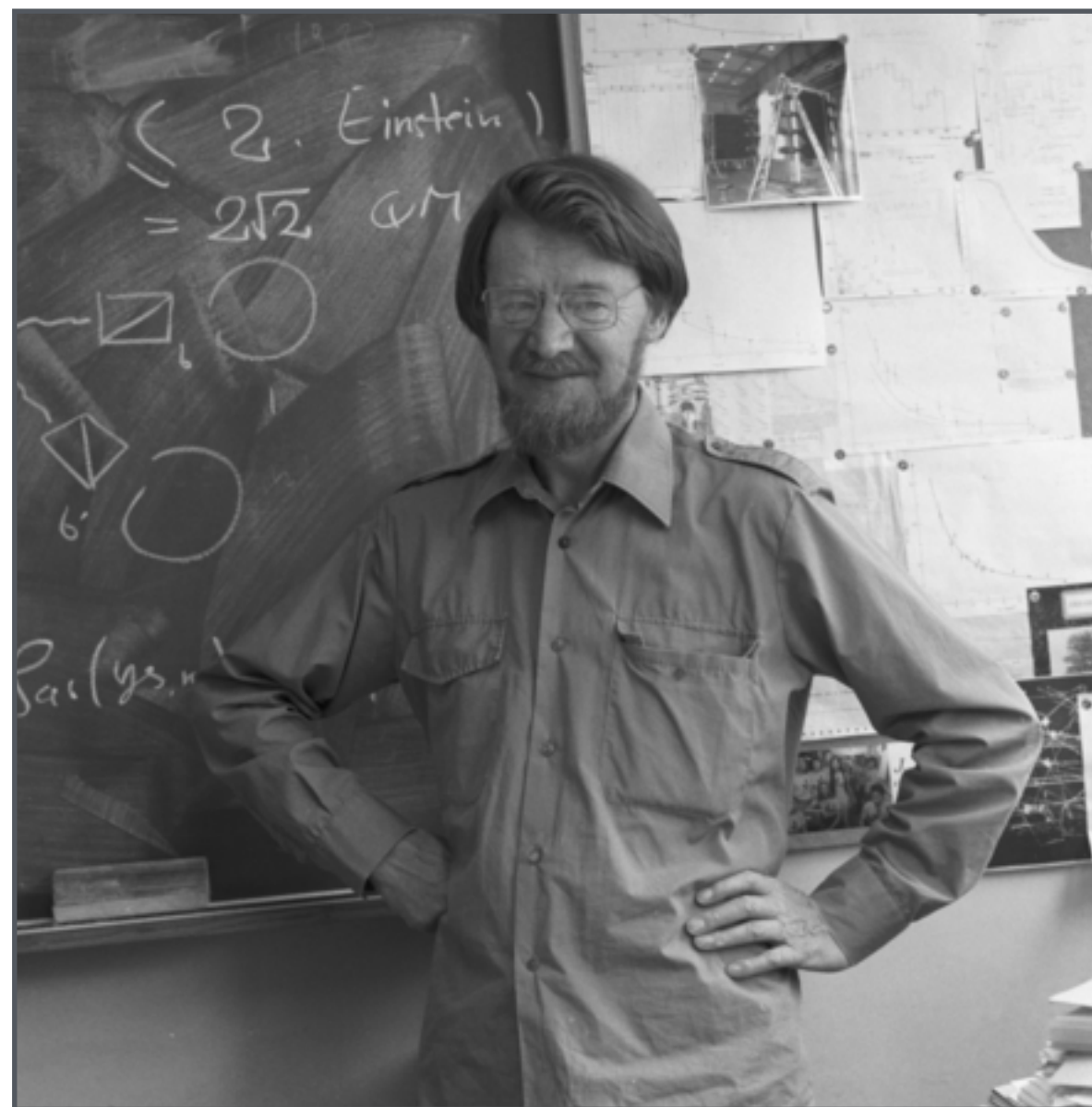
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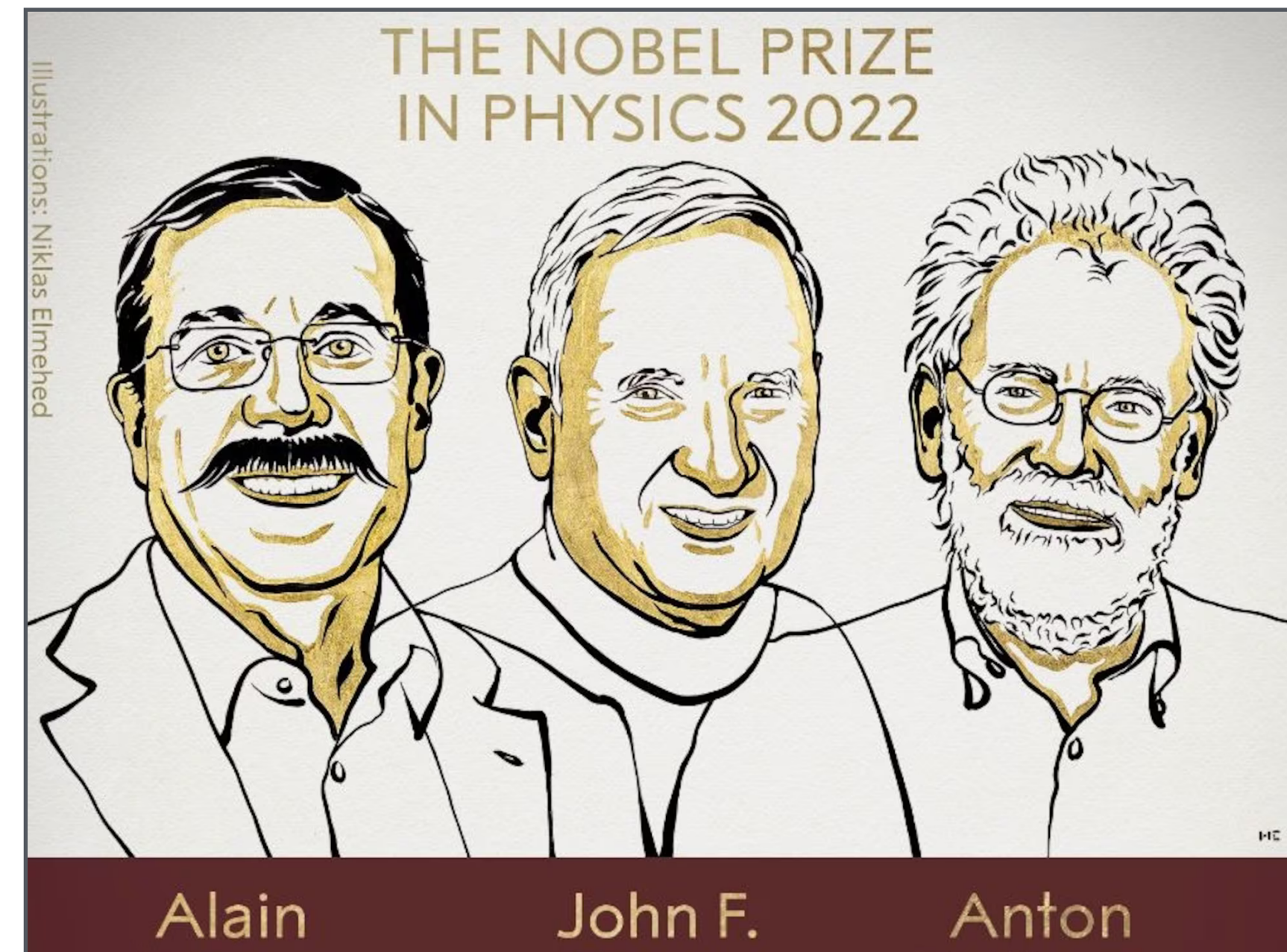
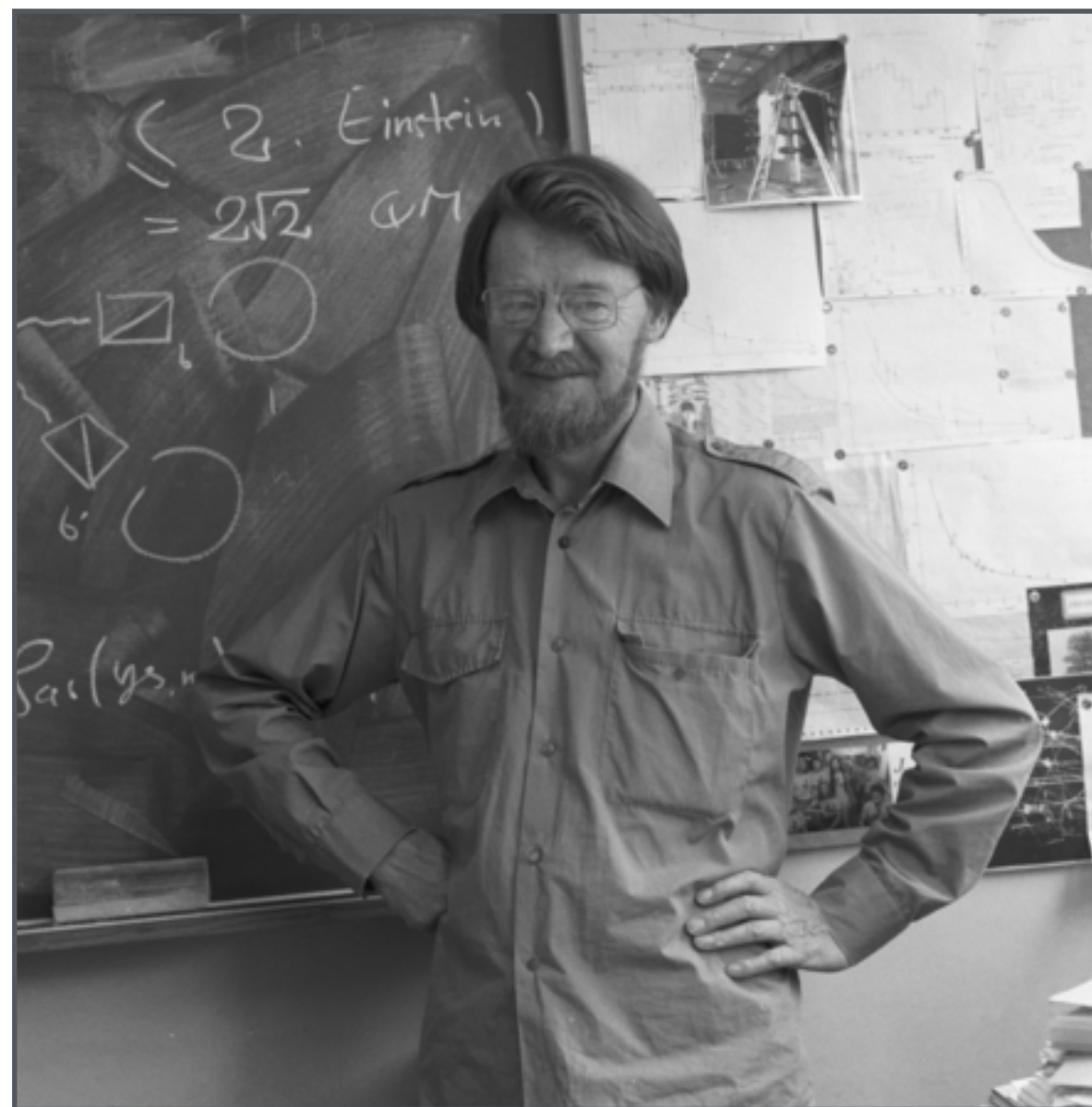
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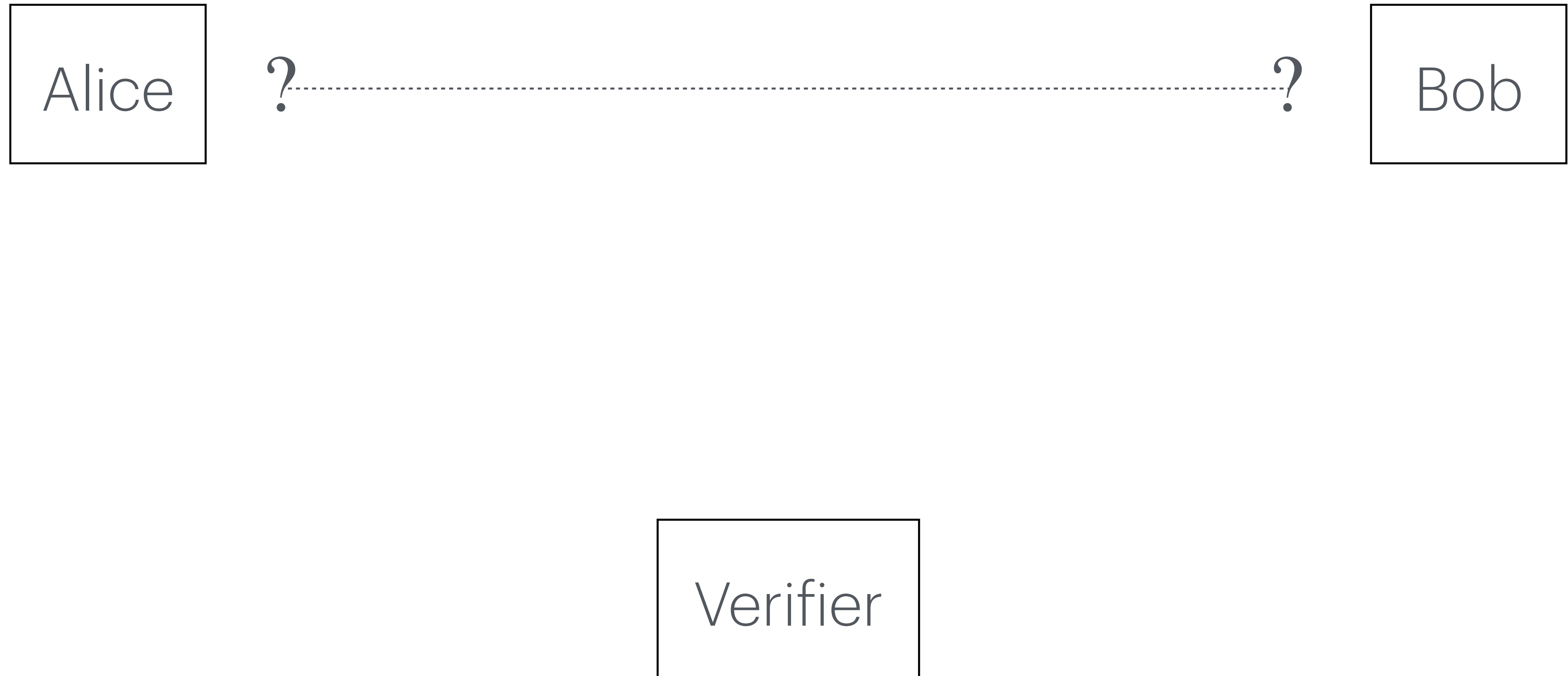


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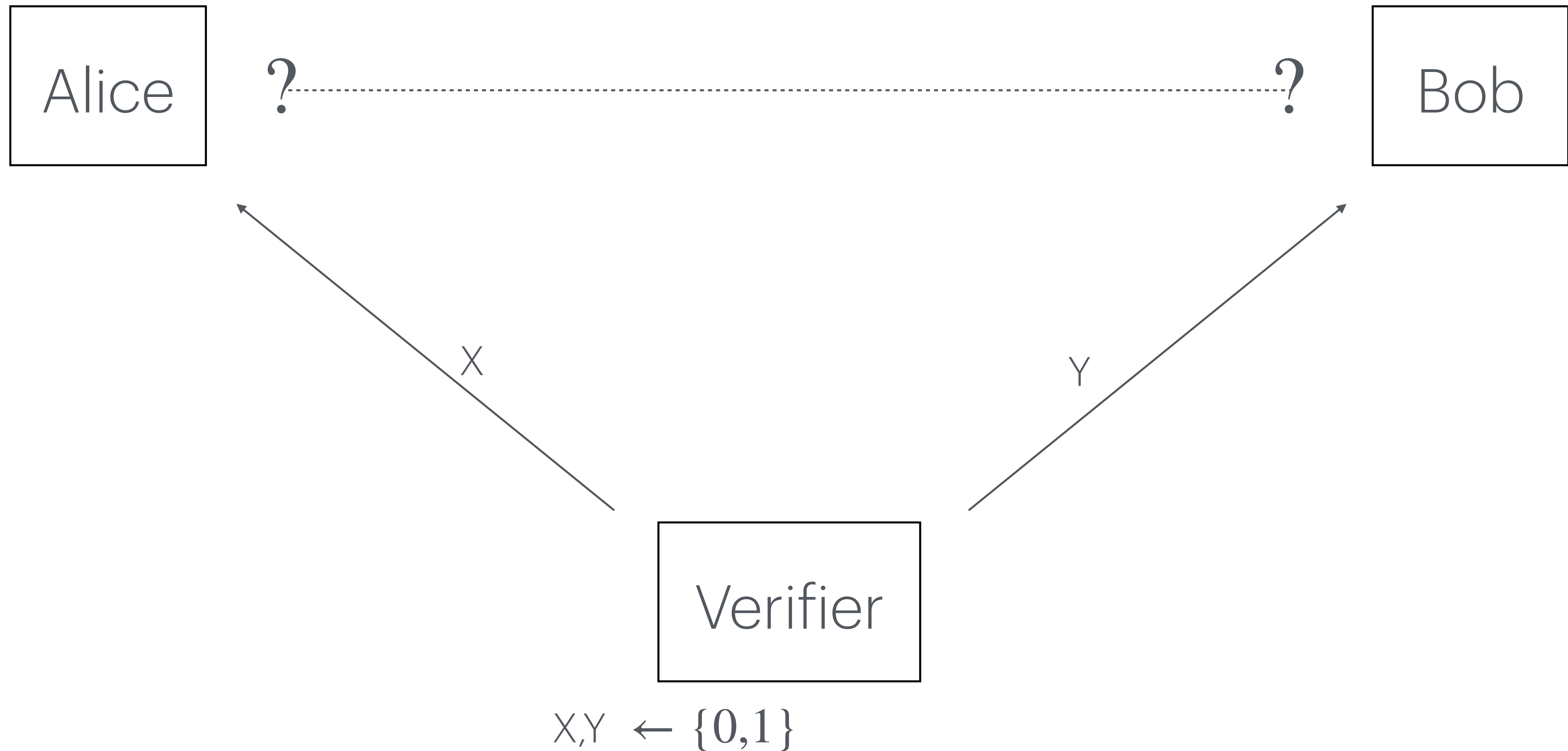


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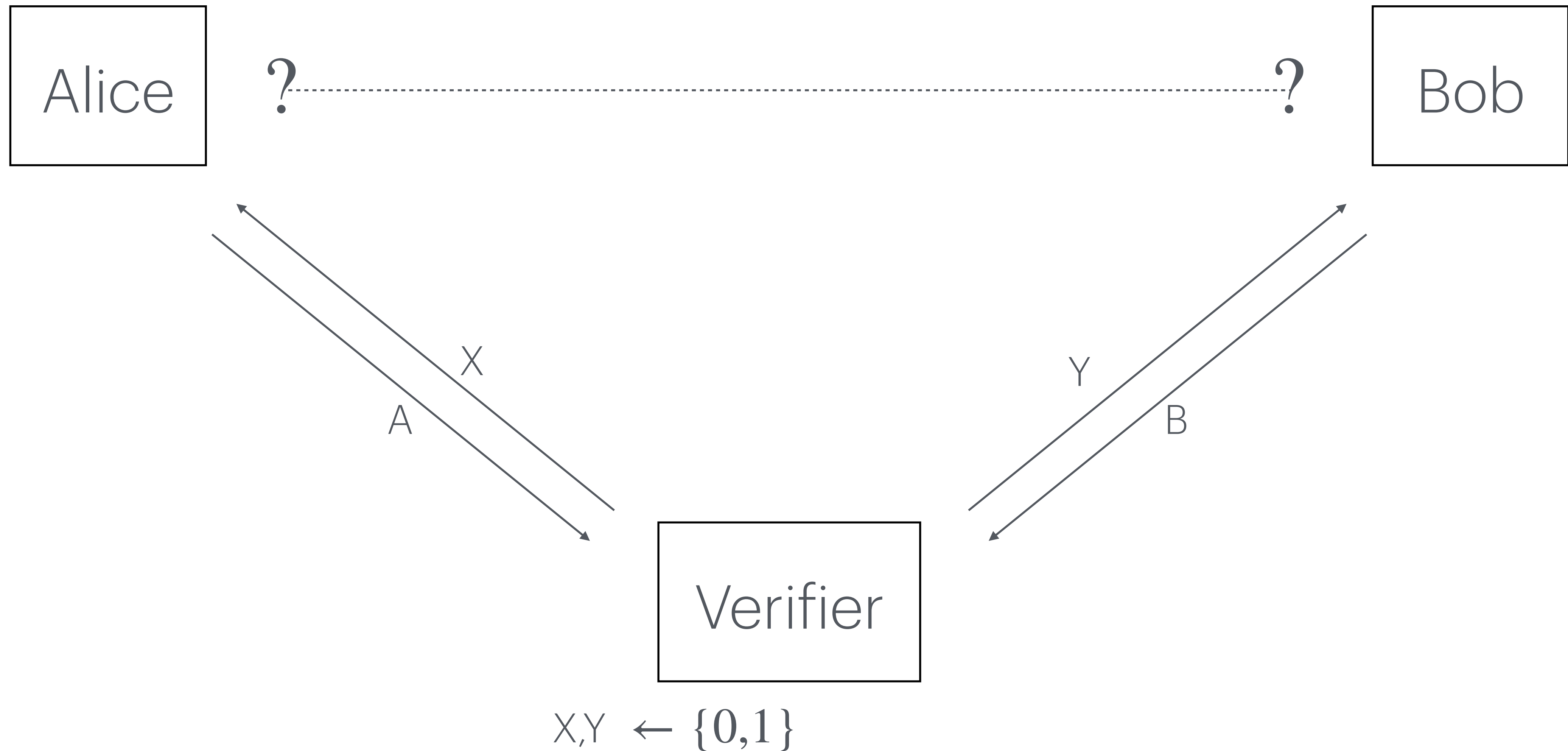




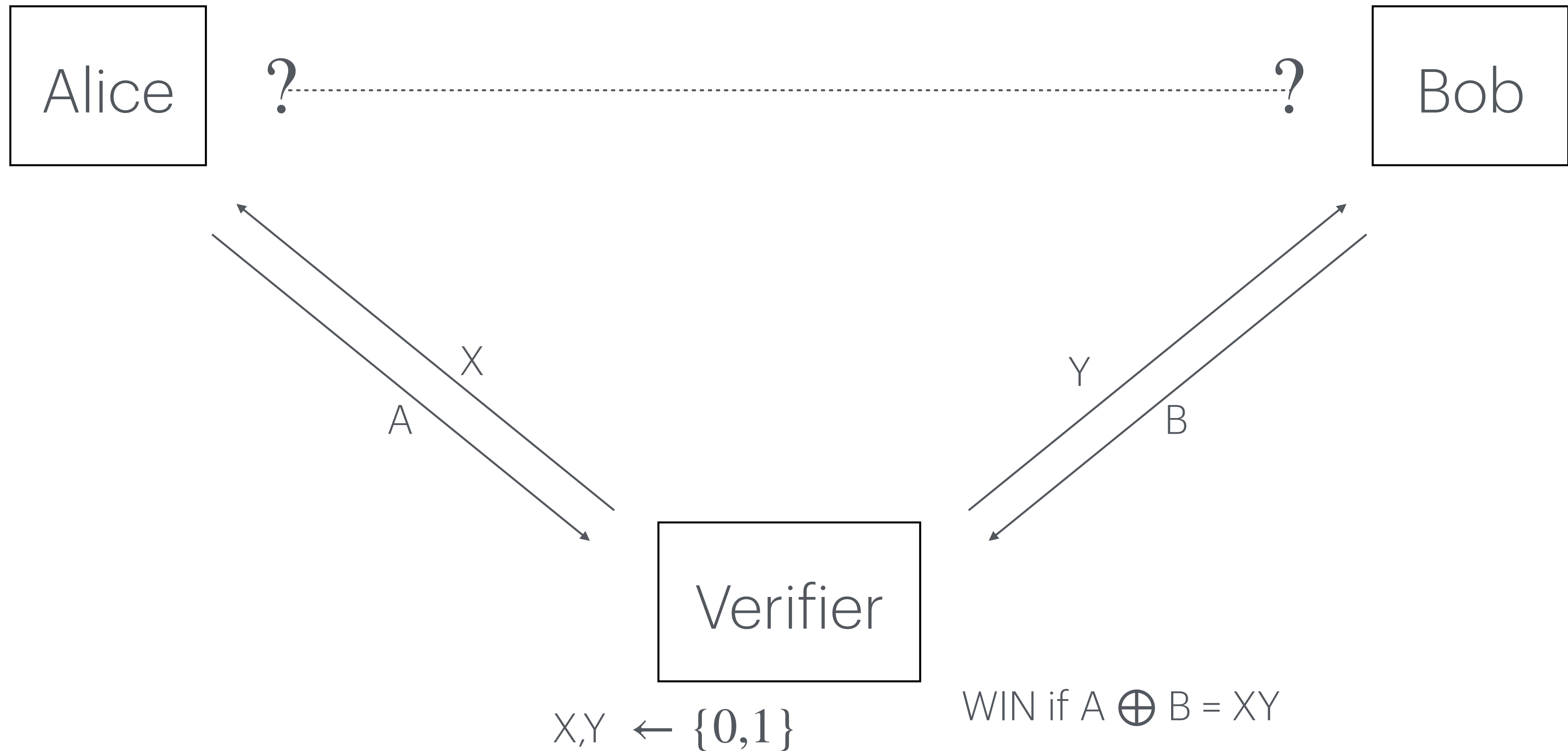
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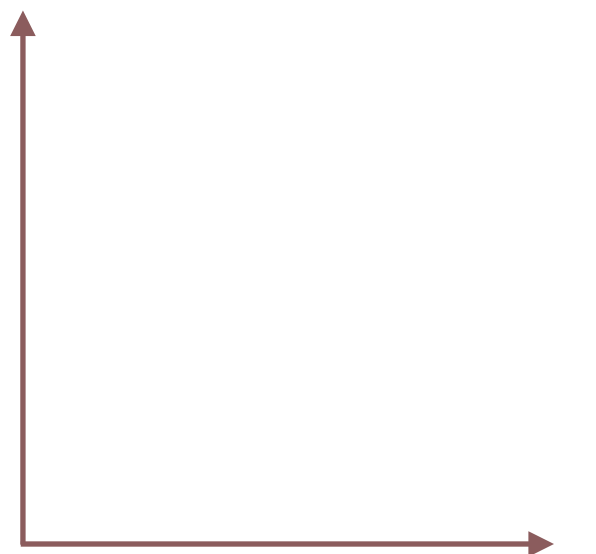


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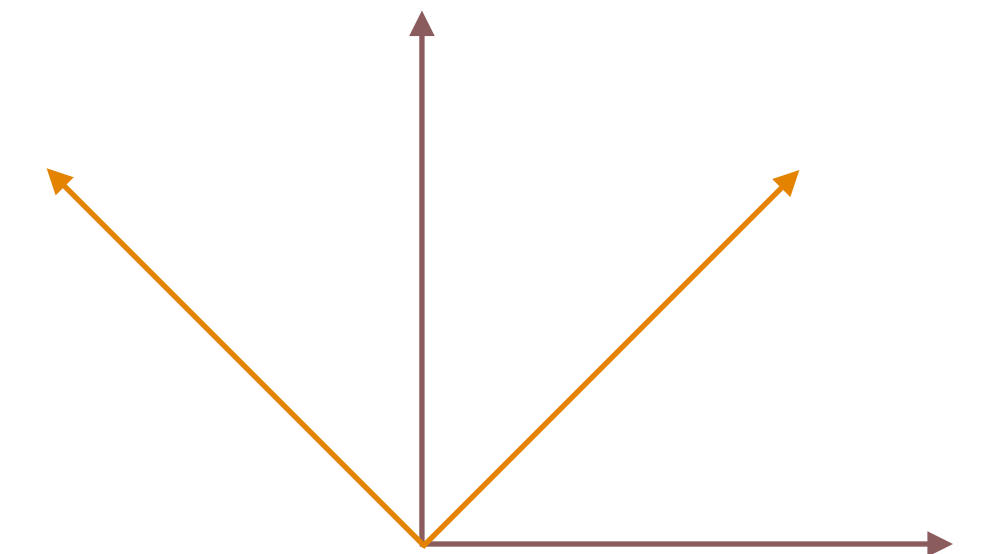
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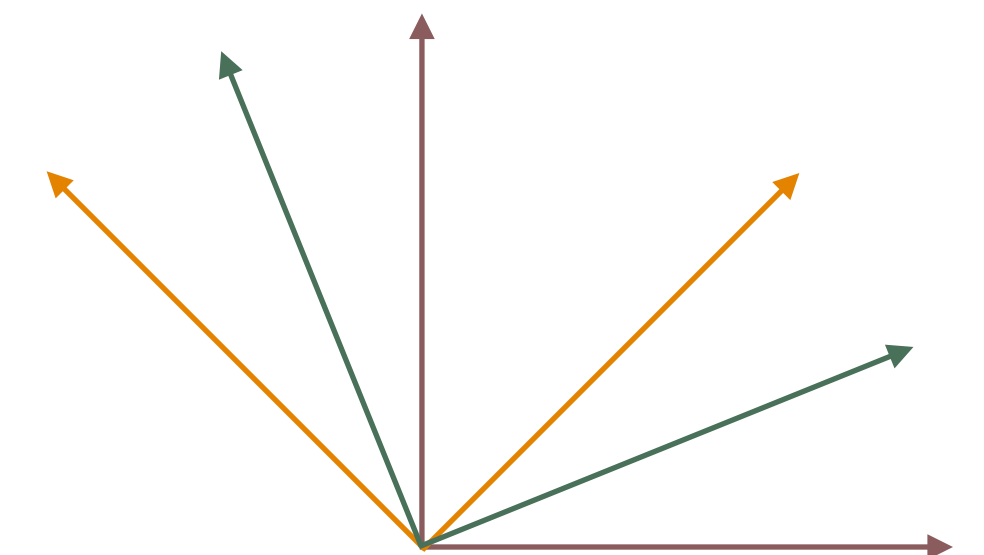
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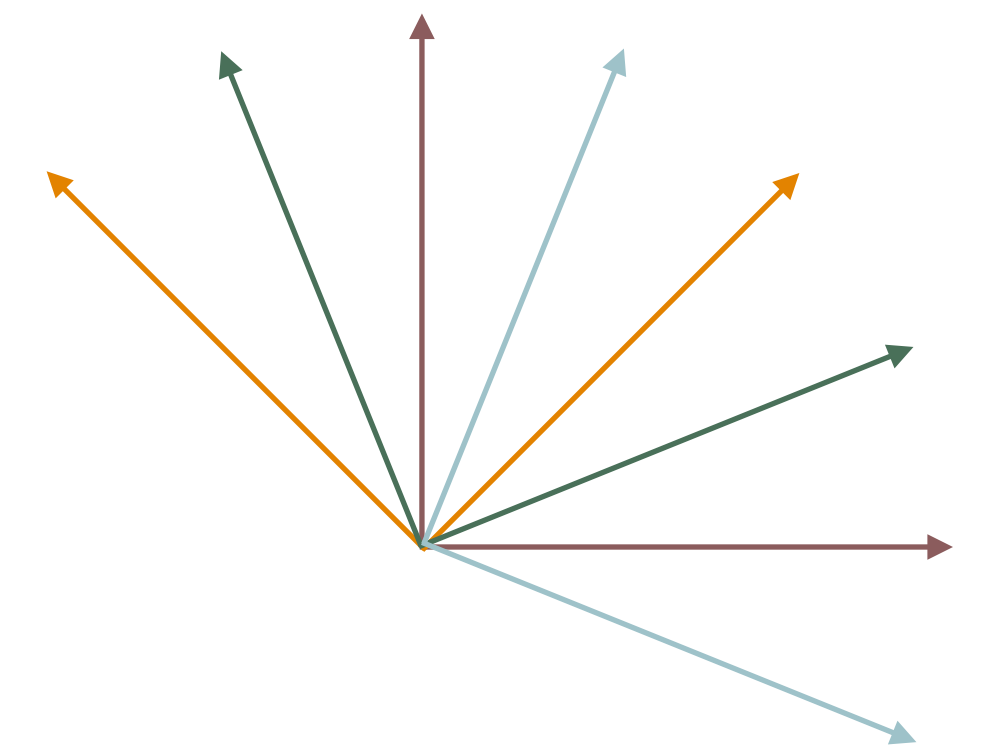
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Thank you!